

Elastic-Degenerate Strings in Bioinformatics: Motivation and Open Problems

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Motivation

- ▶ **Explosion of biological data:** Modern sequencing technologies produce vast amounts of data.
- ▶ **From single sequence to collections:** To ask more complex questions.
- ▶ **Research directions:** Take advantages from analyzing a larger number of highly similar sequences (pangenomics, phylogenomics).
- ▶ **Population Reference Genome:** A concept representing multiple genomes in a unified data structure — enabling more realistic analysis of biological variability.

Approaches

Approach	Idea / Representation	Tools
k-mer-based	Compact colored de Bruijn Graphs k-mer hash table indices	Bifrost, deBGA, Themisto, Cortex, Metagraph, Pufferfish, minimap2
Linearization / Concatenation	Concatenate all sequences: seq1#seq2#... and build global FM-index, r-index	MONI, SPUMONI, MOVI
Variant-aware, MSA-based	Shared regions and differences (SNPs, INDELs) in variation graphs, variable texts	vg, GCSA2, gramtools, GraphAligner, BWBBLE

String

- ▶ **Alphabet** Σ is a finite set of symbols.

DNA: $\Sigma = \{A, C, G, T\}$.

- ▶ A **string** X over Σ is a finite sequence of symbols from Σ :

$$X = s_1 s_2 \dots s_n \in \Sigma^n$$

where $|X| = n$ is the **length** of the string.

- ▶ s_i is the i -th character of X ; $X[i..j]$ is the substring from i to j .
- ▶ ε denotes the **empty string**, i.e., a string of length 0.
- ▶ Special symbols as separators or endmarks #, \$

+ Degeneracy

- ▶ A **degenerate string** X over Σ is a finite sequence of symbols

$$X = S_1 S_2 \dots S_n \in \Sigma^n$$

where S_i is a i -th set of symbols from Σ , i.e., $S_i \subseteq \Sigma$

- ▶ $|X| = n$ is the **length** of the degenerate string.
- ▶ $||X|| = N$ is the **size** of the degenerate string, $N = \sum_{i=0}^n |S_i|$
- ▶ If $|S_i| > 1$, then symbol is **degenerate**.
- ▶ **Language** $L(X)$: set of all strings represented by X

Example:

$$X = \text{GT} \left\{ \begin{array}{c} \text{C} \\ \text{G} \end{array} \right\} \text{AT} \left\{ \begin{array}{c} \text{A} \\ \text{C} \\ \text{T} \end{array} \right\} \text{T}$$

+ Elasticity

- An **Elastic-degenerate string** (EDS)

$$X = S_1 S_2 \dots S_n \in \Sigma^n$$

where S_i is a i -th set of strings over Σ , i.e., $S_i^{i'}$ is a i' -th string in S_i

- **Elasticity:** Every string in $S_i, 0 \leq i \leq n$ may vary in length.
- $|X| = n$ is the **length** of the degenerate string.
- $||X|| = N$ is the **size** of the degenerate string, $N = \sum_{i=0}^n \sum_{i'}^{|S_i|} |S_i^{i'}|$
- **Cardinality** $c = \sum_{i=0}^n |S_i|$

Example:

$$X = \{\text{AAC}\} \left\{ \begin{array}{c} \varepsilon \\ \text{TA} \end{array} \right\} \left\{ \begin{array}{c} \text{CGA} \\ \text{G} \end{array} \right\} \left\{ \begin{array}{c} \text{A} \\ \text{GG} \\ \text{T} \end{array} \right\} \left\{ \begin{array}{c} \varepsilon \\ \text{AA} \end{array} \right\},$$

$$n = 5, N = 15, c = 10$$

$$L(X) = \{\text{AACCGAA}, \text{AACCGAAAA}, \text{AACCGAGG}, \dots, \text{AACTAGT}, \text{AACTAGTAA}\}$$

Pros and Cons

Pros	Cons
Straightforward, easy-to-understand representation	No haplotype paths or variant relationships encoded by default
Can model SNPs and small indels compactly	Cannot represent complex structural variants (e.g. inversions, long inserts)
Polynomial-time algorithms for some tasks	Many problems are NP-hard
Linear construction	Compromise in degeneracy level while construction?

Problem Hardness on EDS

Köppl and Olbrich 2024	Longest Repeating Factor	$\mathcal{O}(N^2)$
	Minimal Unique Substring	NP-hard ¹
	Minimal Absent Word	NP-hard ¹
	Anti-Power Detection	NP-hard ¹
	Longest-Previous-Factor	NP-hard ¹

¹Already hard for binary degenerate strings

Online pattern matching

Elastic-Degenerate String Matching (EDSM) Problem:

Input: EDS X of length n , string pattern P of length m .

Output: All positions in X where at least one occurrence of P ends.

Online Exact EDSM	Time	Space	Impl.
Iliopoulos, Kundu, and S. Pissis 2016	$\mathcal{O}(N + \alpha\gamma nm)$	$\mathcal{O}(m)$	✓
Grossi et al. 2017	$\mathcal{O}(nm^2 + N); \mathcal{O}(N \lceil m/w \rceil)$	$\mathcal{O}(m)$	✓
Aoyama et al. 2018	$\mathcal{O}(nm^{1.5} + N)$	$\mathcal{O}(m)$	✗
Bernardini, Gawrychowski, et al. 2021	$\mathcal{O}(nm^{\omega-1} + N)^1; \mathcal{O}(nm^{1.373} + N)$	Not discussed	✗
S. P. Pissis and Retha 2018	$\mathcal{O}(N \lceil M/w \rceil)$	$\mathcal{O}(M \lceil M/w \rceil)^2$	✓
Cisłak and Grabowski 2020	$\mathcal{O}(occ * m(N/n + \lceil r/w \rceil))^3$	Not discussed	✓

¹ $\omega < 2.373$

² $M = \sum |P_i|$

³ w : word machine size, d : avg number of variants, r : #samples

Online pattern matching

Elastic-Degenerate String Matching (EDSM) with k Mismatches/Errors:

Input: EDS X of length n and cardinality c , string pattern P of length m .

Output: All positions in X where at least one approximate occurrence (with at most k mismatches/errors) of P ends.

Online Aproximate EDSM	Time	Space	Impl.
Bernardini, Pisanti, et al. 2020	$\mathcal{O}(k^2mc + kN)^1; \mathcal{O}(kmc + kN)^2$	$\mathcal{O}(m)$	✗
S. Pissis, Radoszewski, and Zuba 2025	$\mathcal{O}(kmc + kN)^1; \mathcal{O}(k^{2/3}mc + \sqrt{k}N)^2$	$\mathcal{O}(m)$	?
Gawrychowski et al. 2025	$\mathcal{O}(nm^{1.5} + N)$	Not discussed	✗

¹Under Edit distance (EDSM with k Errors)

²Under Hamming distance (EDSM with k Mismatches)

EDS Indexes

	Description	Impl.
Maciuca et al. 2016	BWT FM-index over PRG with anchors	✓
Gibney 2020	$\mathcal{O}(n^\alpha m^\beta)$, $\alpha < 0, \beta < 0$; GST $\mathcal{O}(nm^2)$	✗
Cioni, Guerrini, and Rosone 2024	EDS-BWT based on BWT for sequence collections EBWT	✓

Mapping and Alignment

	Description	Impl.
Mwaniki and Pisanti 2022	Online $\mathcal{O}(mN)$ -time	✗
Büchler, Olbrich, and Ohlebusch 2023	Minimizer based index over PRG	✓

EDS Comparison

Elastic-Degenerate String Intersection (EDSI) problem:

- ▶ **Input:** Two ED-strings X_1 (length n_1 , cardinality c_1 , size N_1) and X_2 (length n_2 , cardinality c_2 , size N_2)
- ▶ **Output:** YES if $L(X_1) \cap L(X_2) \neq \emptyset$, NO otherwise

Reference / Method	Build Time	Lower bound	Impl.
Gabory et al. 2024	$\mathcal{O}(N_1 m_2 + N_2 m_1)$	$\mathcal{O}((N_1 m_2 + N_2 m_1)^{1-\epsilon})^3$	✓

³ $\epsilon > 0$ constant from matrix multiplication lower bounds

Take-Away Message & Summary

- ▶ **Elastic-Degenerate Strings (EDS):**
 - ▶ They can compactly represent both SNPs and indels.
 - ▶ Enable adaptations of string algorithms to variation-aware data.
 - ▶ Solutions exist for some tasks (e.g., pattern matching).
 - ▶ Only a few implementations are haplotype-aware (using phased EDS).
- ▶ How can EDS be built and limited to reflect the best biological image?
- ▶ Is phasing EDS that necessary?
- ▶ What level of degeneracy is common in real data - in bacteria, eukaryota?
- ▶ Do you see other application?

Very nice review on EDS and other variable strings in Ascone et al. 2024.

Thank you for your attention

supervisor: prof. Ing. Jan Holub, Ph.D.

ROBOPROX 

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SGS23/205/OHK3/3T/18

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